### Scheduling

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#### Goals of Process Scheduling

- Maximize CPU utilization
- Fair CPU allocation
- Minimize Turnaround Time
- Minimize Waiting Time
- Maximize the number of processes that complete execution per time unit ("throughput")

#### **Process Scheduling Times**

#### Arrival Time

The time at which the process arrives in the ready queue.

#### Completion Time

The time at which a process is completed.

#### Burst Time

CPU time that a process requires for its execution.

#### Turnaround Time

= Completion Time – Arrival Time

#### Waiting Time

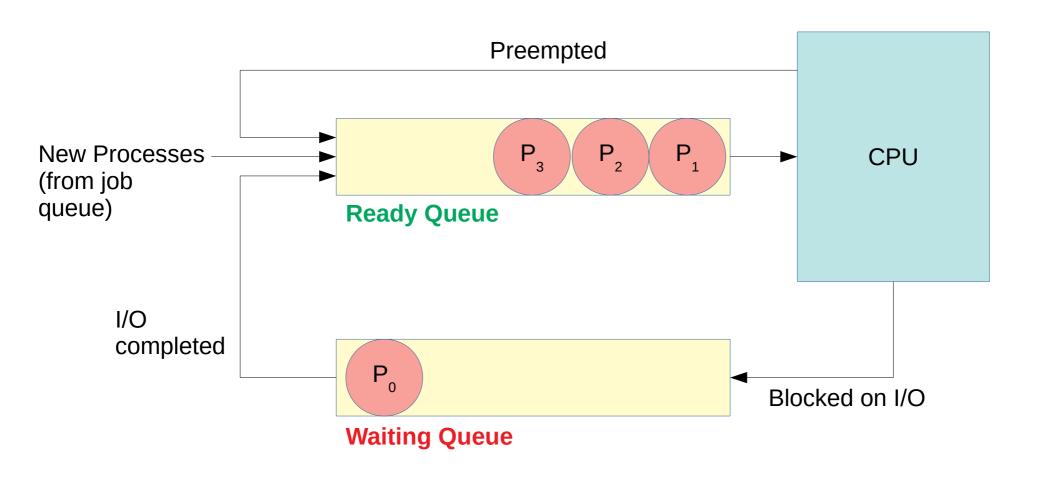
The amount of time a process has been waiting in the ready queue.



#### Two Types of Scheduling Algorithms

- Non-preemptive Scheduling
   A process holds the CPU until it terminates or it switches from running to waiting state.
- Preemptive Scheduling
   A running process can be taken away from the CPU in favor of other processes.

### Scheduling Queues (simplified)





#### Scheduling Algorithms

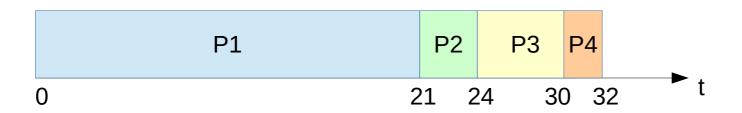
- First Come First Serve
- Shortest Job First
- Shortest Remaining Time First
- Round Robin Scheduling
- Priority Based Scheduling
- Highest Response Ratio Next
- Multilevel Queue Scheduling
- Multilevel Feedback Queue Scheduling



#### First Come First Serve (FCFS) (1)

| Process | <b>Burst Time</b> |
|---------|-------------------|
| P1      | 21                |
| P2      | 3                 |
| P3      | 6                 |
| P4      | 2                 |

Assumption: Arrival Time of all processes is 0



#### First Come First Serve (FCFS) (2)

- FIFO (First-In-First-Out) queue structure
- Simple to implement
- Non-preemptive
- Does not consider priority
- Low throughput possible due to convoy effect (Long processes will delay execution of short processes)
- No starvation (assuming that every process will eventually complete)



### Shortest Job First (SJF) (1)

| Process | <b>Burst Time</b> |
|---------|-------------------|
| P1      | 21                |
| P2      | 3                 |
| P3      | 6                 |
| P4      | 2                 |

Assumption: Arrival Time of all processes is 0

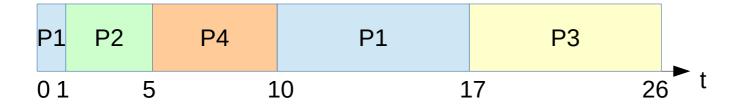
| P4  | P2  | P3   | P1 |                 |
|-----|-----|------|----|-----------------|
| 0 2 | 2 5 | 5 11 | 32 | $\frac{1}{2}$ t |

### Shortest Job First (SJF) (2)

- Always selects the process with the smallest burst time for execution first.
- A process' priority is the inverse of its predicted CPU burst time.
- The burst time of each process must be known in advance.
- Risk of starvation.

## Shortest Remaining Time First (SRTF) (1)

| Process | <b>Burst Time</b> | <b>Arrival Time</b> |
|---------|-------------------|---------------------|
| P1      | 8                 | 0                   |
| P2      | 4                 | 1                   |
| P3      | 9                 | 2                   |
| P4      | 5                 | 3                   |



# Shortest Remaining Time First (SRTF) (2)

- Whenever a process with a shorter burst time arrives, the currently executed process is preempted.
- Overhead due to context switching
- Starvation is possible (if short processes are continually added)

## Highest Response Ratio Next (HRRN)

- Similar to SJN with a small modification
- Decision which job is next is based on the highest response ratio.

$$response\ ratio = 1 + \frac{waiting\ time}{estimated\ run\ time}$$

- Jobs that have spent a relatively long waiting time will be preferred
- Mitigates the problem of starvation

### Round Robin Scheduling (RR) (1)

| Process | <b>Burst Time</b> |
|---------|-------------------|
| P1      | 21                |
| P2      | 3                 |
| P3      | 6                 |
| P4      | 2                 |

Quantum: 5

Assumption: Arrival Time of all processes is 0

|   | P1 | P2  | P3 | P4   | P1 | P3   | P1  | P1 | P1 |    |   |
|---|----|-----|----|------|----|------|-----|----|----|----|---|
| 0 | 5  | 5 8 | 3  | 13 1 | .5 | 20 2 | 1 2 | 26 | 31 | 32 | t |

### Round Robing Scheduling (RR) (2)

- Assigns a fixed time quantum per process and cycles through all processes
- Process is rescheduled if it does not complete within the quantum
- No starvation
- Overhead due to context switching, especially with small time quanta
  - Quantum should be significantly higher than context switch time, e.g. 100 ms when context switch time is < 10  $\mu$ s

#### Exercise

- Determine the average waiting time for the previous examples of the following scheduling algorithms:
  - First Come First Serve (FCFS)
  - Shortest Job First (SJF)
  - Shortest Remaining Time First (SRTF)
  - Round Robin (RR)

### Priority Based Scheduling (1)

| Process | <b>Burst Time</b> | Priority |
|---------|-------------------|----------|
| P1      | 21                | 2        |
| P2      | 3                 | 1        |
| P3      | 6                 | 4        |
| P4      | 2                 | 3        |

Assumption: Arrival Time of all processes is 0



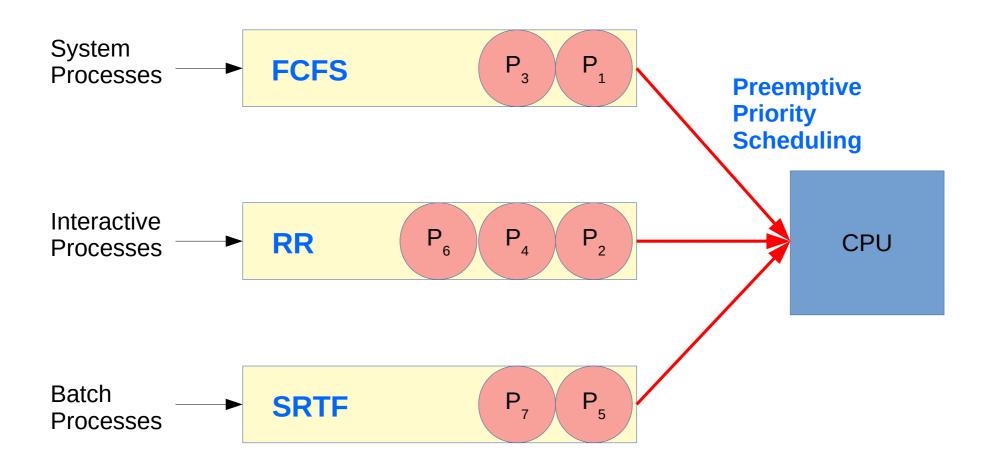
### Priority Based Scheduling (2)

- Process with the highest priority is executed first
- Processes with the same priority require a second scheduling algorithm, e.g. FCFS
- Risk of starvation in preemptive priority scheduling

### Multilevel Queue Scheduling (1)

- Processes are classified into different groups each of which has its own scheduling requirements.
- Ready queue consists of multiple separate queues each of which has its own scheduling algorithm.
- Additional scheduling is required among the queues (e.g. preemptive priority scheduling)

# Multilevel Queue Scheduling (Example)

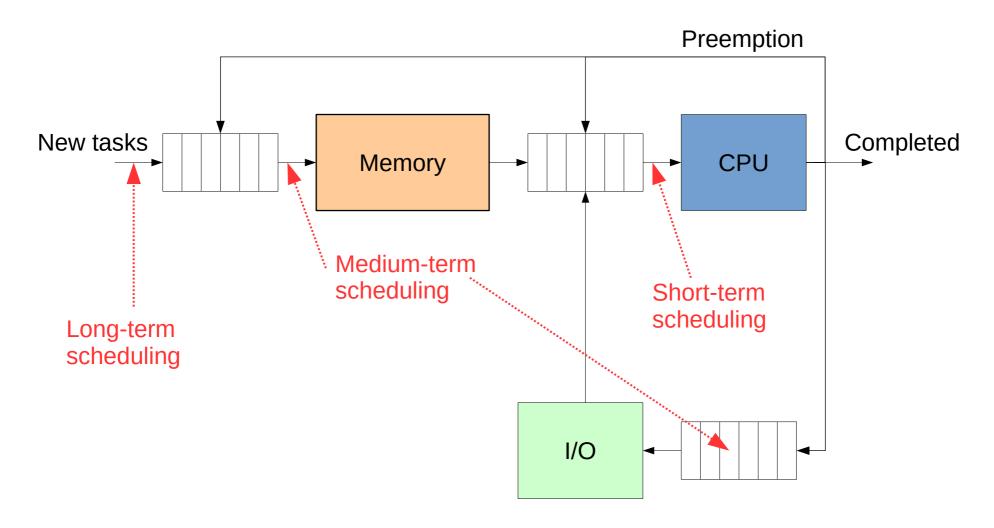


### Multilevel Feedback Queue Scheduling

- Enhanced variant of Multilevel Queue Scheduling
- Allows processes to move between the queues based on their spent CPU time
- Processes in high priority queues that spend too much CPU time may be moved to a lower priority queue
- Processes that have been waiting too long in a lower priority queue may be moved to a higher priority queue
- Most general scheduler, but also most complex to implement



### Scheduling Strategies in the OS



#### Lab Exercises (1)

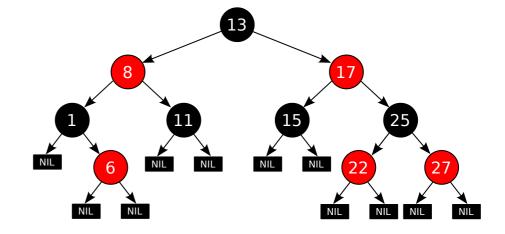
- Implement simulations of the following scheduling algorithms in C
  - First Come First Serve (FCFS)
  - Shortest Remaining Time First (SRTF)
  - Round Robin Scheduling (RR)

#### Lab Exercises (2)

- What concept does Linux provide from a user's perspective to run processes with different priorities?
- Set up a scenario in which processes with higher priority are granted more CPU time than others.
- What mechanisms does Linux provide to prevent users from depriving other users of CPU time?

## Completely Fair Scheduler (CFS) (1)

- Presently used by the Linux Kernel (since 2.6.23, October 2007)
- Developed by Ingo Molnár
- Idea: Ideal Fairness
  - Every Process receives 1/n
     CPU time
  - Sleeping processes "earn"
     CPU time, i.e. they are given a boost when they wake up

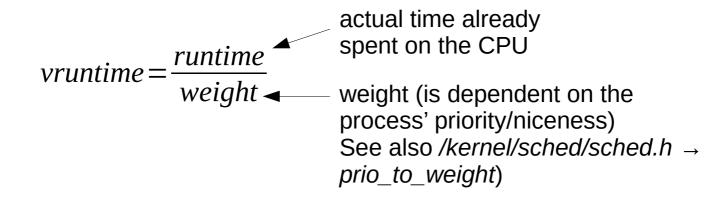


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- CFS uses a red-black tree instead of queues as its underlying data structure
- CFS uses nanoseconds granularity accounting

# Completely Fair Scheduler (CFS) (2)

Each process has a virtual runtime (vruntime)



- The higher the weight, the lower the impact of the actual runtime. → Slower increase of virtual runtime for processes with higher priority.
- CFS always chooses the process with the lowest virtual runtime.



## Completely Fair Scheduler (CFS) (3)

- Some related files and tools
  - /proc/stat
  - /sys/kernel/debug/sched\_features
  - proc/sys/kernel/sched\_\*
  - taskset
  - trace-cmd
  - kernelshark
- More technical details on CFS
  - https://www.kernel.org/doc/html/latest/scheduler/sched-design-CF S.html
  - https://doc.opensuse.org/documentation/leap/tuning/html/book-sletuning/cha-tuning-taskscheduler.html



# Real-time Operating Systems (RTOS)

- Guarantees a certain behavior within predictable (well-defined) time constraints.
- OS-internal actions must be finished
- Two groups of RTOS
  - Hard
    - Deadlines must always be met
    - Example: Collision avoidance systems in aircrafts
  - Soft
    - Occasionally missed deadlines do not pose critical risks
    - Example: Decoders for media streaming



#### Earliest Deadline First (EDF)

- Process with the earliest deadline gets dispatched to the CPU
- Algorithm requires a system clock (absolute time)
- Algorithm is preemptive
- All deadlines can be met, as long as sufficient computing resources are available
- Disadvantages
  - In overload scenarios, deadlines will be missed unpredictably
  - Difficult to implement in hardware



#### **Exercises**

• Develop the timing diagram for EDF scheduling of two processes  $P_1$ ,  $P_2$ , and  $P_3$  with the following parameters:

| Process | Arrival<br>Time | Burst<br>Time | Deadline |
|---------|-----------------|---------------|----------|
| P1      | 0               | 4             | 20       |
| P2      | 1               | 5             | 15       |
| P3      | 2               | 6             | 16       |

- Rate-Monontonic Scheduling (RMS) is a fixed-priority realtime scheduling algorithm that is usually preferred over EDF.
  - What are its advantages over EDF?
  - How does it work?

#### Questions for Review

- What are the goals of process scheduling?
- What is a major disadvantage of SJF and SRTF?
- Why should the time quantum in RR scheduling be significantly higher than the time required to perform a context switch?
- Explain the principle of how a Multilevel Feedback Queue Scheduler works.
- What is the definition of an RTOS?
- What happens in EDF scheduling when the CPU load exceeds 100%?